The Framework of Linear Nested Sequents

Björn Lellmann

Vienna, Nov. 4, 2015

Sequent systems and modal logics

Sequent calculi for modal logics are well-established and well-understood – but not entirely satisfactory!

Some desiderata for "good" calculi [Wansing:'02]:

- subformula property: all the material in the premisses is contained in the conclusion
- separation: distinct left and right introduction rules
- locality: no restrictions on the context
- modularity: obtain other logics by changing single rules, e.g., following Došen's Principle: only vary structural rules

Sequent systems and modal logics

It can be easily verified that each of the standard rule systems [for modal logics] fails to satisfy some of the philosophical requirements [...].

[Wansing:'94]

E.g.:

	$\dfrac{\Gamma \Rightarrow A}{\Box\Gamma \Rightarrow \Box A}\ \mathsf{k}$	$\dfrac{\Box\Gamma \Rightarrow A}{\Box\Gamma \Rightarrow \Box A}\ 4$ $\dfrac{\Gamma, A \Rightarrow \Delta}{\Gamma, \Box A \Rightarrow \Delta}\ \mathsf{t}$
Subformula property:	✓	✓
Separation:	✗	✓
Locality:	✗	✗
Modularity:	✗	

Solutions: structures with sequents in them

The solution according to internal approaches:

Extend the sequent structure!

By now, there are many ways to do so:

- Higher-level sequents : Sequents of sequents of sequents of...
  [Došen:'85]
- 2-sequents: Streams of sequents
  [Masini:'92]
- Display calculi: structural connectives for all operators
  [Belnap:'82, Wansing:'94, Kracht:'96]
- Nested sequents: Trees of sequents
  [Kashima:'94, Brünnler:'06, Poggiolesi:'09]
- ...

The Question

What is the simplest extension of the sequent structure satisfying these desiderata for modal logics?

Reminder: Modal logics

The formulae of modal logic are given by

$$\varphi ::= \mathsf{Var} \mid \neg\varphi \mid \varphi \wedge \varphi \mid \varphi \vee \varphi \mid \varphi \to \varphi \mid \Box\varphi$$

The Hilbert-style presentation of normal modal logic K is given by the axioms for classical propositional logic and

$$\mathsf{k}\quad \Box(A \to B) \wedge \Box A \to \Box B \qquad\qquad \frac{\vdash A}{\vdash \Box A}\ \mathsf{nec}$$

The standard sequent system contains the standard propositional rules together with

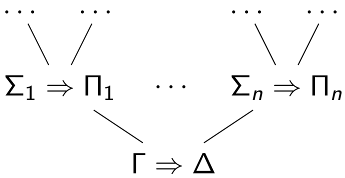
$$\frac{\Gamma \Rightarrow A}{\Box\Gamma \Rightarrow \Box A}\ \mathsf{k}$$

Case study: Nested sequents

Definition ([Brünnler:'09,Poggiolesi:'09])

A nested sequent is a finite tree whose nodes are labelled with sequents. The interpretation ι of this nested sequent is

$$\bigwedge\Gamma \rightarrow \bigvee\Delta \vee \bigvee_{i=1}^{n} \Box\iota(\Sigma_i \Rightarrow \Pi_i) \; .$$



Fact

The nested sequent calculus with modal rules $\Box_R$ and $\Box_L$ is sound and cut-free complete for modal logic K.

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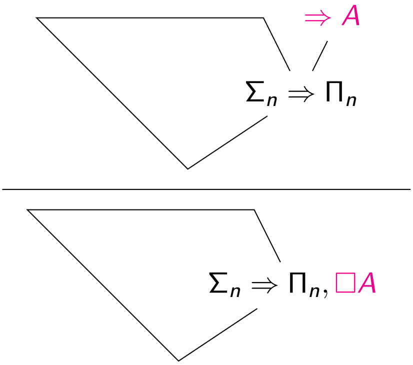

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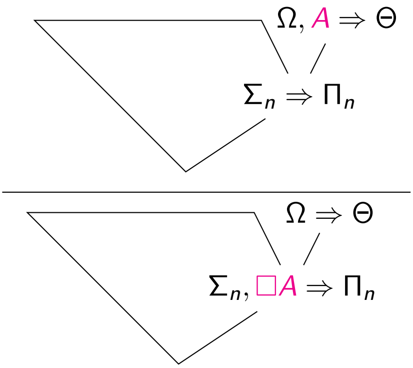

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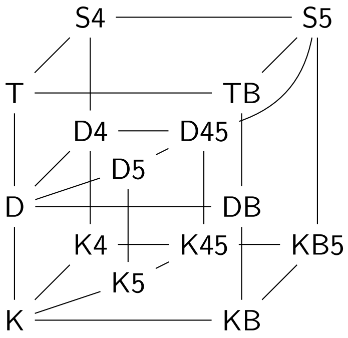
Case study: Nested sequents

In fact, there are cut-free modular nested sequent systems for all logics in the (normal) modal cube:



[Marin, Straßburger:'14]

Case study: Nested sequents

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	seq.	nested seq.
Subformula property:	✓	✓
Separation:	✗	✓
Locality:	✗	✓
Modularity:	✗	✓

Trees are nice, but can we go simpler?

A different approach: 2-sequents

$$\begin{array}{ccc} \vdots & & \vdots \\ \varepsilon & & \varepsilon \\ \Gamma_n & \Rightarrow & \Delta_n \\ \vdots & & \vdots \\ \Gamma_1 & & \Delta_1 \end{array}$$

Definition ([Masini:'92])

A 2-sequent is an infinite, eventually empty stream of sequents. It's interpretation is

$$\bigwedge\Gamma_1 \to \bigvee\Delta_1 \vee \Box(\dots\Box(\bigwedge\Gamma_n \to \bigvee\Delta_n)\dots)$$

Fact

The 2-sequent calculus with modal rules $\Box_R$ and $\Box_L$ is sound and cut-free complete for modal logic KD.

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$$\frac{\begin{array}{ccc} \vdots & & \varepsilon \\ \varepsilon & & A \\ \Gamma_n & \Rightarrow & \Delta_n \\ \vdots & & \vdots \\ \Gamma_1 & & \Delta_1 \end{array}}{\begin{array}{ccc} \vdots & & \vdots \\ \varepsilon & & \varepsilon \\ \Gamma_n & \Rightarrow & \Delta_n, \Box A \\ \vdots & & \vdots \\ \Gamma_1 & & \Delta_1 \end{array}}$$

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$$\dfrac{\begin{array}{ccc} \vdots & & \vdots \\ \Sigma, A & & \Pi \\ \Gamma_n & \Rightarrow & \Delta_n \\ \vdots & & \vdots \\ \Gamma_1 & & \Delta_1 \end{array}}{\begin{array}{ccc} \vdots & & \vdots \\ \Sigma & & \Pi \\ \Gamma_n, \Box A & \Rightarrow & \Delta_n \\ \vdots & & \vdots \\ \Gamma_1 & & \Delta_1 \end{array}}$$

A different approach: 2-sequents

This sequent system for KD *admits cut-elimination, [...] and the introduction rules are separate, symmetrical and explicit, but no indication is given of how to present axiomatic extensions [...]. [I]t is not clear how Masini's framework may be modified in order to obtain a 2-sequent calculus for* K*.*

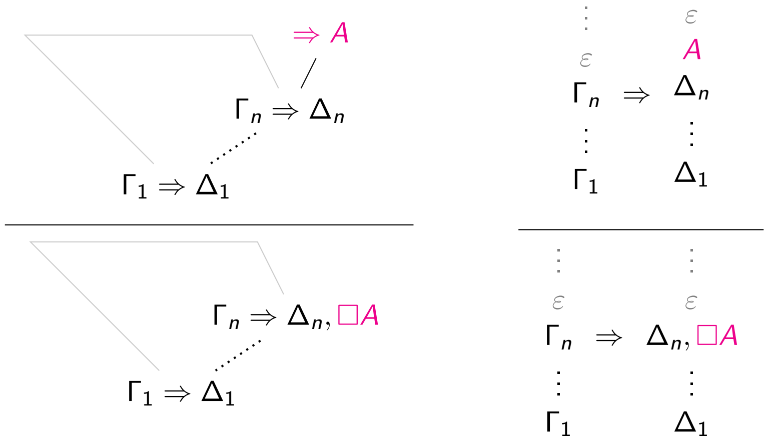
[Wansing:'02]

	seq.	nested seq.	2-seq.
Subformula property:	✓	✓	✓
Separation:	✗	✓	✓
Locality:	✗	✓	✓
Modularity:	✗	✓	✗

Infinite linear structures are nice, but can we go simpler?

Connections between nested sequents and 2-sequents

Comparing the rules reveals they are essentially the same, e.g.:



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$$
\frac{\begin{array}{c} \Sigma, A \Rightarrow \Pi \\ / \\ \Gamma_n \Rightarrow \Delta_n \\ \vdots \\ \Gamma_1 \Rightarrow \Delta_1 \end{array}}{\begin{array}{c} \Sigma \Rightarrow \Pi \\ / \\ \Gamma_n, \Box A \Rightarrow \Delta_n \\ \vdots \\ \Gamma_1 \Rightarrow \Delta_1 \end{array}}
\qquad
\frac{\begin{array}{ccc} \vdots & & \vdots \\ \Sigma, A & & \Pi \\ \Gamma_n & \Rightarrow & \Delta_n \\ \vdots & & \vdots \\ \Gamma_1 & & \Delta_1 \end{array}}{\begin{array}{ccc} \vdots & & \vdots \\ \Sigma & & \Pi \\ \Gamma_n, \Box A & \Rightarrow & \Delta_n \\ \vdots & & \vdots \\ \Gamma_1 & & \Delta_1 \end{array}}
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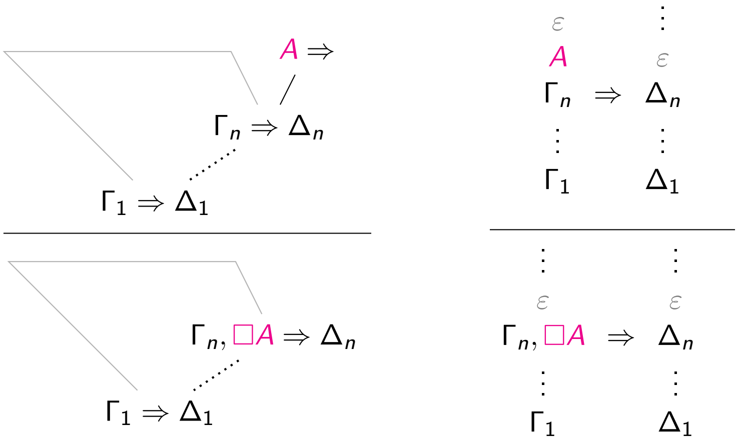


So the structure of finite lists of sequents is enough for KD!

Let's try finite lists of sequents!

Linear nested sequents

Definition

A linear nested sequent (LNS) is a finite list of sequents, written

$$\Gamma_1 \Rightarrow \Delta_1 /\!\!/ \dots /\!\!/ \Gamma_n \Rightarrow \Delta_n$$

and interpreted as $\wedge\Gamma_1 \rightarrow \vee\Delta_1 \vee \Box(\dots\Box(\wedge\Gamma_n \rightarrow \vee\Delta_n)\dots)$.

The nested sequent system for K yields the modal rules of $\mathsf{LNS_K}$:

$$\frac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta /\!\!/ \Sigma, A \Rightarrow \Pi /\!\!/ \mathcal{H}}{\mathcal{G} /\!\!/ \Gamma, \Box A \Rightarrow \Delta /\!\!/ \Sigma \Rightarrow \Pi /\!\!/ \mathcal{H}}\ \Box_L \qquad \frac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta /\!\!/ \ \Rightarrow A}{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, \Box A}\ \Box_R$$

The propositional rules are standard, e.g.:

$$\frac{\mathcal{G} /\!\!/ \Gamma, B \Rightarrow \Delta /\!\!/ \mathcal{H} \quad \mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, A /\!\!/ \mathcal{H}}{\mathcal{G} /\!\!/ \Gamma, A \rightarrow B \Rightarrow \Delta /\!\!/ \mathcal{H}}\ \rightarrow_L \qquad \frac{\mathcal{G} /\!\!/ \Gamma, A \Rightarrow \Delta, B /\!\!/ \mathcal{H}}{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, A \rightarrow B /\!\!/ \mathcal{H}}\ \rightarrow_R$$

Completeness for linear nested sequents

We could show completeness via cut elimination ... but it's easier!

Observation: The data structure of LNS is the same as that of a history in backwards proof search for a sequent calculus.

So we simply simulate a sequent derivation in the last components: ($\mathcal{G}$ is the history)

$$\dfrac{\dfrac{\Gamma \Rightarrow A}{\Box\Gamma \Rightarrow \Box A}\ \mathsf{k}}{\vdots\ \mathcal{G}} \quad \rightsquigarrow \quad \dfrac{\dfrac{\mathcal{G} /\!\!/ \Box\Gamma \Rightarrow \Box A /\!\!/ \Gamma \Rightarrow A}{\mathcal{G} /\!\!/ \Box\Gamma \Rightarrow \Box A /\!\!/ \Rightarrow A}\ \Box_L}{\mathcal{G} /\!\!/ \Box\Gamma \Rightarrow \Box A}\ \Box_R$$

Theorem
LNS_K *is sound and cut-free complete for* K.

Corollary: Cut-free completeness of the nested sequent calculus.

Extensions

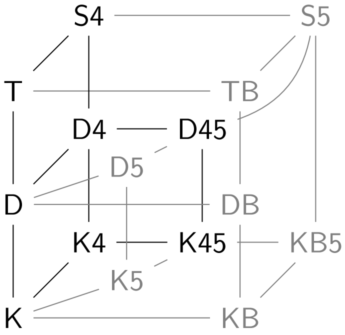
Extensions, e.g. (lifted shamelessly from nested sequent calculi):

$$\frac{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta /\!/ A \Rightarrow}{\mathcal{G} /\!/ \Gamma, \Box A \Rightarrow \Delta}\ \mathsf{d} \qquad \frac{\mathcal{G} /\!/ \Gamma, A \Rightarrow \Delta /\!/ \mathcal{H}}{\mathcal{G} /\!/ \Gamma, \Box A \Rightarrow \Delta /\!/ \mathcal{H}}\ \mathsf{t}$$

$$\frac{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta /\!/ \Sigma, \Box A \Rightarrow \Pi /\!/ \mathcal{H}}{\mathcal{G} /\!/ \Gamma, \Box A \Rightarrow \Delta /\!/ \Sigma \Rightarrow \Pi /\!/ \mathcal{H}}\ 4 \qquad \frac{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta /\!/ \Sigma \Rightarrow \Pi, \Box A /\!/ \mathcal{H}}{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta, \Box A /\!/ \Sigma \Rightarrow \Pi /\!/ \mathcal{H}}\ 5$$

Theorem

The LNS calculi for extensions of K *with axioms from* $\mathsf{d}, \mathsf{t}, 4$ *or* $\mathsf{d}, 4, (4 \wedge 5)$ *are cut-free complete and modular.*



Application: intuitionistic logic

The same idea connects Maehara's multi-succedent calculus and Fitting's nested sequent calculus for first-order intuitionistic logic:

The intuitionistic interpretation of $\Gamma_1 \Rightarrow \Delta_1 /\!\!/ \dots /\!\!/ \Gamma_n \Rightarrow \Delta_n$ is

$$\bigwedge\Gamma_1 \to \bigvee\Delta_1 \vee (\bigwedge\Gamma_2 \to \bigvee\Delta_2 \vee (\dots(\bigwedge\Gamma_n \to \bigvee\Delta_n)\dots))$$

Restricting the nested sequent rules yields the rules of $\mathsf{LNS}_{\mathsf{Int}}$, e.g.:

$$\frac{\mathcal{G} /\!\!/ \Gamma, B \Rightarrow \Delta /\!\!/ \mathcal{H} \quad \mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, A /\!\!/ \mathcal{H}}{\mathcal{G} /\!\!/ \Gamma, A \to B \Rightarrow \Delta /\!\!/ \mathcal{H}} \to_L \qquad \frac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta /\!\!/ A \Rightarrow B}{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, A \to B} \to_R$$

$$\frac{\mathcal{G} /\!\!/ \Gamma, A\alpha \Rightarrow \Delta /\!\!/ \mathcal{H}}{\mathcal{G} /\!\!/ \Gamma, \forall x.Ax \Rightarrow \Delta /\!\!/ \mathcal{H}} \forall_L \qquad \frac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta /\!\!/ \ \Rightarrow A\alpha}{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, \forall x.Ax} \forall_R$$

if α in $\mathcal{H}$ then α in $\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta$ $\qquad$ α not in the conclusion

$$\frac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta /\!\!/ \Sigma, A \Rightarrow \Pi /\!\!/ \mathcal{H}}{\mathcal{G} /\!\!/ \Gamma, A \Rightarrow \Delta /\!\!/ \Sigma \Rightarrow \Pi /\!\!/ \mathcal{H}} \text{Lift}$$

(Rules for $\exists$ analogous and other rules local.)

Intuitionistic logic: Completeness

We simulate Maehara's rules in the last components, e.g.:

$$\dfrac{\dfrac{\Gamma, A \Rightarrow B}{\Gamma \Rightarrow \Delta, A \to B}\ {\to_R}}{\vdots\ \mathcal{G}} \quad \rightsquigarrow \quad \dfrac{\dfrac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, A \to B /\!\!/ \Gamma, A \Rightarrow B}{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, A \to B /\!\!/ A \Rightarrow B}\ \text{Lift}}{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, A \to B}\ {\to_R}$$

$$\dfrac{\dfrac{\Gamma \Rightarrow A\alpha}{\Gamma \Rightarrow \Delta, \forall x.Ax}\ \forall_R}{\vdots\ \mathcal{G}} \quad \rightsquigarrow \quad \dfrac{\dfrac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, \forall x.Ax /\!\!/ \Gamma \Rightarrow A\alpha}{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, \forall x.Ax /\!\!/ \Rightarrow A\alpha}\ \text{Lift}}{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, \forall x.Ax}\ \forall_R$$

The other rules are easy.

Theorem

$\mathsf{LNS}_{\mathsf{Int}}$ *is sound and complete for first-order intuitionistic logic.*

Corollary: Cut-free completeness of Fitting's calculus.

Taking it further: non-normal modal logics

The language of monotone modal logic M is that of modal logic.

The Hilbert-style presentation of M is given by axioms for classical propositional logic and the rule

$$\frac{\vdash A \to B}{\vdash \Box A \to \Box B}\ \text{Mon}$$

The sequent system for M contains the standard propositional rules and the modal rule

$$\frac{A \Rightarrow B}{\Box A \Rightarrow \Box B}\ \text{Mon}$$

Common extensions:

MCN = K

MC $\Box A \land \Box B \to \Box(A \land B)$ MN $\Box\top$

M

Non-normal linear nested sequents

To capture the sequent rule $\dfrac{A \Rightarrow B}{\Box A \Rightarrow \Box B}$ we use a marker $\mathbin{\backslash\!\!\backslash}$ for "unfinished rules": A non-normal LNS has the form $(n \geq 0)$

$$\Gamma_1 \Rightarrow \Delta_1 /\!\!/ \dots /\!\!/ \Gamma_n \Rightarrow \Delta_n$$
$$\text{or} \quad \Gamma_1 \Rightarrow \Delta_1 /\!\!/ \dots /\!\!/ \Gamma_n \Rightarrow \Delta_n \mathbin{\backslash\!\!\backslash} \Gamma_{n+1} \Rightarrow \Delta_{n+1}$$

Translating the sequent rule Mon yields the modal rules of $\mathsf{LNS_M}$:

$$\frac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta \mathbin{\backslash\!\!\backslash} \ \Rightarrow B}{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta, \Box B}\ \Box_R \qquad \frac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta /\!\!/ \Sigma, A \Rightarrow \Pi}{\mathcal{G} /\!\!/ \Gamma, \Box A \Rightarrow \Delta \mathbin{\backslash\!\!\backslash} \Sigma \Rightarrow \Pi}\ \Box_L$$

The propositional rules cannot be applied inside $\mathbin{\backslash\!\!\backslash}$.

The completeness proof for $\mathsf{LNS_M}$ then uses the simulation

$$\begin{array}{c} \dfrac{A \Rightarrow B}{\Box A \Rightarrow \Box B}\ \text{Mon} \\ \vdots\ \mathcal{G} \end{array} \quad \rightsquigarrow \quad \dfrac{\dfrac{\mathcal{G} /\!\!/ \Box A \Rightarrow \Box B /\!\!/ A \Rightarrow B}{\mathcal{G} /\!\!/ \Box A \Rightarrow \Box B \mathbin{\backslash\!\!\backslash} \ \Rightarrow B}\ \Box_L}{\mathcal{G} /\!\!/ \Box A \Rightarrow \Box B}\ \Box_R$$

Modularity for non-normal modal logics

C $\Box A \land \Box B \to \Box(A \land B)$

$$\frac{\mathcal{G} \backslash\!\backslash \Gamma \Rightarrow \Delta}{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta}\ \mathsf{c}$$

N $\Box\top$

$$\frac{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta}{\mathcal{G} \backslash\!\backslash \Gamma \Rightarrow \Delta}\ \mathsf{n}$$

D_1 $\neg\Box\bot$

$$\frac{\mathcal{G} \backslash\!\backslash \ \Rightarrow}{\mathcal{G}}\ \mathsf{d}_1$$

D_2 $\neg(\Box A \land \Box\neg A)$

$$\frac{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta \backslash\!\backslash A \Rightarrow}{\mathcal{G} /\!/ \Gamma, \Box A \Rightarrow \Delta}\ \mathsf{d}_2$$

T $\Box A \to A$

$$\frac{\mathcal{G} /\!/ \Gamma, A \Rightarrow \Delta}{\mathcal{G} /\!/ \Gamma, \Box A \Rightarrow \Delta}\ \mathsf{t}$$

4 $\Box A \to \Box\Box A$

$$\frac{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta /\!/ \Sigma, \Box A \Rightarrow \Pi}{\mathcal{G} /\!/ \Gamma, \Box A \Rightarrow \Delta \backslash\!\backslash \Sigma \Rightarrow \Pi}\ 4$$

5 $\Box\neg A \lor \Box\neg\Box A$

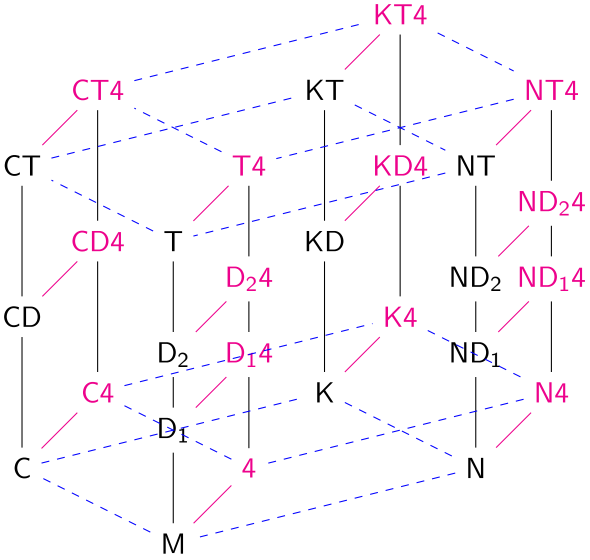
$$\frac{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta /\!/ \Sigma \Rightarrow \Pi, \Box A}{\mathcal{G} /\!/ \Gamma \Rightarrow \Delta, \Box A \backslash\!\backslash \Sigma \Rightarrow \Pi}\ 5$$

Theorem

For $\mathcal{A} \subseteq \{\mathsf{C}, \mathsf{N}, \mathsf{D}_1, \mathsf{D}_2, \mathsf{T}, 4\}$ the calculus $\mathsf{LNS}_{\mathsf{M}\mathcal{A}}$ is sound and complete for $\mathsf{M}\mathcal{A}$. Similar for some combinations with 5.

(Use calculi e.g. from [Lavendhomme-Lucas:'00, Indrzejczak:'05].)

Modularity: the modal tesseract



(Restoring the bridge between normal and non-normal logics)

The desiderata

	seq.	nested seq.	2-seq	LNS
Subformula property:	✓	✓	✓	✓
Separation:	✗	✓	✓	✓
Locality:	✗	✓	✓	✓
Modularity:	✗	✓	✗	✓

What about connections to other frameworks?

Hypersequents

The data structure of LNS is rather familiar from another setting:

Definition ([Avron:'96])

A hypersequent is a finite list of sequents, written

$$\Gamma_1 \Rightarrow \Delta_1 \mid \cdots \mid \Gamma_n \Rightarrow \Delta_n$$

and interpreted as $\Box(\bigwedge\Gamma_1 \to \bigvee\Delta_1) \vee \cdots \vee \Box(\bigwedge\Gamma_n \to \bigvee\Delta_n)$.

This interpretation suggests the external structural rules:

$$\frac{\mathcal{G} \mid \Sigma \Rightarrow \Pi \mid \Gamma \Rightarrow \Delta \mid \mathcal{H}}{\mathcal{G} \mid \Gamma \Rightarrow \Delta \mid \Sigma \Rightarrow \Pi \mid \mathcal{H}}\ \mathsf{EEX}$$

$$\frac{\mathcal{G} \mid \Gamma \Rightarrow \Delta \mid \Gamma \Rightarrow \Delta \mid \mathcal{H}}{\mathcal{G} \mid \Gamma \Rightarrow \Delta \mid \mathcal{H}}\ \mathsf{EC} \qquad \frac{\mathcal{G} \mid \mathcal{H}}{\mathcal{G} \mid \Gamma \Rightarrow \Delta \mid \mathcal{H}}\ \mathsf{EW}$$

They are part of almost all hypersequent calculi for modal logics.

Hypersequents and linear nested sequents

Observation 1: EC and EW are the structural nested sequent rules for (4) and (t) (modulo internal structural rules):

$$\frac{\mathcal{G} \mid \Gamma \Rightarrow \Delta \mid \Gamma \Rightarrow \Delta \mid \mathcal{H}}{\mathcal{G} \mid \Gamma \Rightarrow \Delta \mid \mathcal{H}}\ \mathsf{EC} \qquad \text{vs.} \qquad \frac{\mathcal{G} /\!\!/ \Gamma \Rightarrow \Delta /\!\!/ \Sigma \Rightarrow \Pi /\!\!/ \mathcal{H}}{\mathcal{G} /\!\!/ \Gamma, \Sigma \Rightarrow \Delta, \Pi /\!\!/ \mathcal{H}}\ \dot{\mathsf{t}}$$

$$\frac{\mathcal{G} \mid \mathcal{H}}{\mathcal{G} \mid \Gamma \Rightarrow \Delta \mid \mathcal{H}}\ \mathsf{EW} \qquad \text{vs.} \qquad \frac{\mathcal{G} /\!\!/ \mathcal{H}}{\mathcal{G} /\!\!/ \Rightarrow /\!\!/ \mathcal{H}}\ \overline{4}$$

Observation 2: $\mathsf{LNS_K} + \dot{\mathsf{t}} + \overline{4} + \mathsf{EEX}$ is (essentially) the hypersequent calculus for S5 from [Restall:'07].

Theorem

$\mathsf{LNS_K} + \dot{\mathsf{t}} + \overline{4} + \mathsf{EEX}$ *is sound and cut-free complete for* S5 *(under the* LNS-*interpretation).*

(Similarly we obtain e.g. Avron's calculus etc.)

Hypersequents and intuitionistic logic

So let's try to add EEX to a different LNS calculus!

Theorem

$\mathsf{LNS}_{\mathsf{Int}} + \mathsf{EEX}$ *is sound and cut-free complete for first-order classical logic.*

This yields calculi for classical and intuitionistic logic satisfying Wansing's desiderata and following Došen's principle.

Question: Can we view EEX as "backtracking" in game semantics?

$$\dfrac{\dfrac{\dfrac{\Rightarrow \bot /\!\!/ A \Rightarrow A}{A \Rightarrow \bot /\!\!/ \ \Rightarrow A}\ \text{Lift}}{\Rightarrow A /\!\!/ A \Rightarrow \bot}\ \text{EEX / backtracking}}{\Rightarrow A, A \to \bot}\ \to_R$$

Conclusion

Summing up:

- Finite lists of sequents give good systems for normal and non-normal modal logics and intuitionistic logic
- An easy method to show cut-free completeness
- A connection to hypersequents via external exchange.

Future work:

- Complexity of proof search (partly done)
- Syntactic cut elimination for LNS
- "Proper" LNS systems for logics without cut-free sequent calculi (e.g., modal logic B, constant-domain Int).